

Exercises to the lecture
Concurrency Theory
Sheet 4

Roland Meyer, Viktor Vafeiadis

Delivery until 20.05.2014 at 12h

Exercise 4.1

Consider the following procedure CAS ("Compare-And-Swap"), where $r = \text{CAS}(x, e, n)$ means that the procedure returns the final value of r .

```

r = CAS(x, e, n):
  atomic (
    r = [x]
    if r = e then ([x] = n; r = 1)
    else(r = 0)
  )

```

- a) Complete the following rule so that it specifies the behaviour of CAS:

$$\frac{\dots}{r = \text{CAS}(x, e, n) \text{ sat } (\boxed{P}, R, G, \boxed{Q})}$$

- b) Implement a lock using the procedure CAS.
Do not use "assume", but you can use "if" and "while".
- c) Prove that your lock implementation satisfies

$$\frac{\begin{array}{l} \text{sem_stable}(P, R) \quad \text{sem_stable}(Q, R) \\ P \Rightarrow p \mapsto - * F \quad p \mapsto \text{TID} * F \Rightarrow Q \\ \{p \mapsto 0 \rightsquigarrow p \mapsto \text{TID}\} \subseteq G \end{array}}{\text{lock}(p, \text{TID}) \text{ sat } (\boxed{P}, R, G, \boxed{Q})}$$

Exercise 4.2

Let $\text{ls}(x, z)$ be defined as $x = z \wedge \mathbf{emp} \vee \exists y. x \mapsto y * \text{ls}(y, z)$.

- a) Prove $(x \mapsto y \text{ } \text{---} \otimes (P * Q)) \Rightarrow ((x \mapsto y \text{ } \text{---} \otimes P) * Q) \vee (P * (x \mapsto y \text{ } \text{---} \otimes Q))$.
- b) Prove $((x \mapsto y) \text{ } \text{---} \otimes \text{ls}(k, m)) \Rightarrow (\text{ls}(k, x) * \text{ls}(y, m))$.
- c) Conclude $\text{sem_stable}(\text{ls}(k, m), R)$ with

$$R := \{x \mapsto y \rightsquigarrow (x \mapsto z * z \mapsto y), (x \mapsto z * z \mapsto y) \rightsquigarrow x \mapsto y\}.$$

Exercise 4.3

Prove the soundness of the PRIM rule:

$$\frac{\vdash_{\text{SL}} \{P\} c \{Q\} \quad \text{writes}(c) \cap \text{fv}(R, G) = \emptyset}{c \text{ sat } (P, R, G, Q)}$$

Delivery until 20.05.2014 at 12h into the box next to 34-401.4