

Exercises to the lecture
Complexity Theory
Sheet 4

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Delivery until 20.11.2017 at 18h

Exercise 4.1 (Alternation Bounded QBF II)

Recall the definition of *alternation bounded* QBF of Exercise 3.2. There, we have shown that $\Sigma_i\text{QBF}$ ($\Pi_i\text{QBF}$) lies in Σ_i^P (Π_i^P). In this exercise, you shall prove that Σ_i ($\Pi_i\text{QBF}$) is Σ_i^P -hard (Π_i^P -hard) with respect to polynomial time reductions.

Hint: Take an arbitrary language in Σ_i^P and reduce it to $\Sigma_i\text{QBF}$. In the lecture we have shown that QBF is PSPACE-complete. Extract the idea from this proof.

Exercise 4.2 (Collapse of the Polynomial Hierarchy)

In this exercise, we prove what is sufficient for a collapse of the polynomial hierarchy.

- a) Assume we have an $i \in \mathbb{N}$ so that $\Sigma_i^P = \Sigma_{i+1}^P$. Show that $\Pi_i^P = \Pi_{i+1}^P$.
- b) Prove the following: If we have an $i \in \mathbb{N}$ with $\Sigma_i^P = \Pi_i^P$, then for any $i' \geq i$ we have that $\Sigma_{i'}^P = \Pi_{i'}^P = \Sigma_i^P$. Hence, the polynomial hierarchy collapses to the i -th level.
Hint: Prove the statement by an induction on i' . For the induction step, make use of the fact that $\Sigma_{i'+1}\text{QBF}$ is $\Sigma_{i'+1}^P$ -hard.
- c) Show that the existence of an $i \in \mathbb{N}$ with $\Sigma_i^P = \Sigma_{i+1}^P$ is already sufficient to cause a collapse of the polynomial hierarchy to the i -th level.

Exercise 4.3 (Minimal Boolean formulas)

Two Boolean formulas are called *equivalent* if they have the same value under any assignment to the variables. A formula φ is called *minimal* if there is no smaller formula that is equivalent to φ . We set $\text{EQUIV} = \{(\varphi, \psi) \mid \varphi \text{ and } \psi \text{ are equivalent}\}$ and $\text{MIN} = \{\varphi \mid \varphi \text{ is minimal}\}$.

- a) Show that EQUIV is in Π_1^P .
- b) Conclude that MIN is a problem in Π_2^P .

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